



# Idealized BGPsec: Formally Verifiable BGP

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for the

Informal BGPsec Design Group

# Informal BGPsec Group

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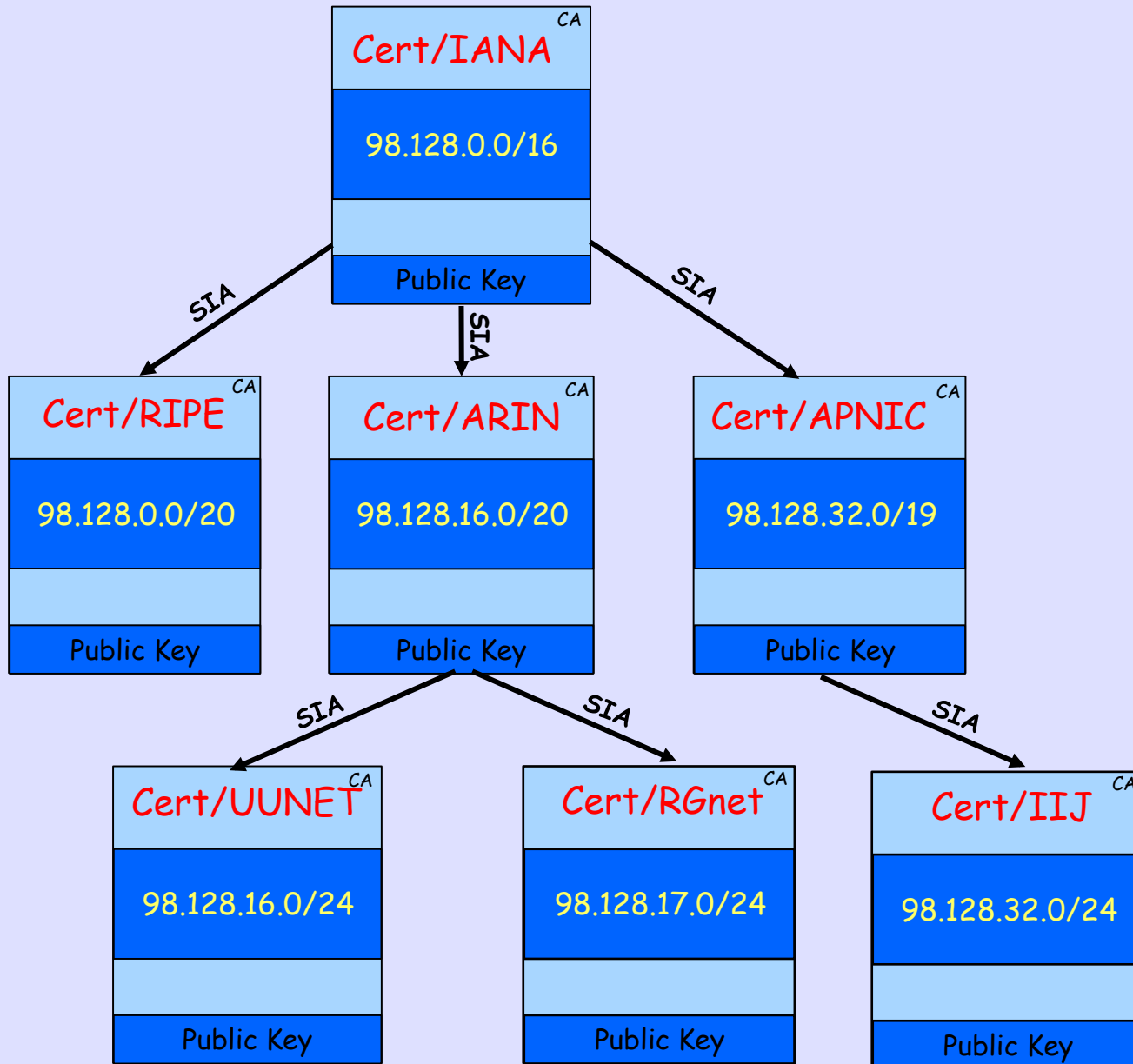
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# Assume RPKI



# Assume RPKI-RTR

**RPKI Portal GUI**

[split](#)  
[roa](#)  
[delete](#)

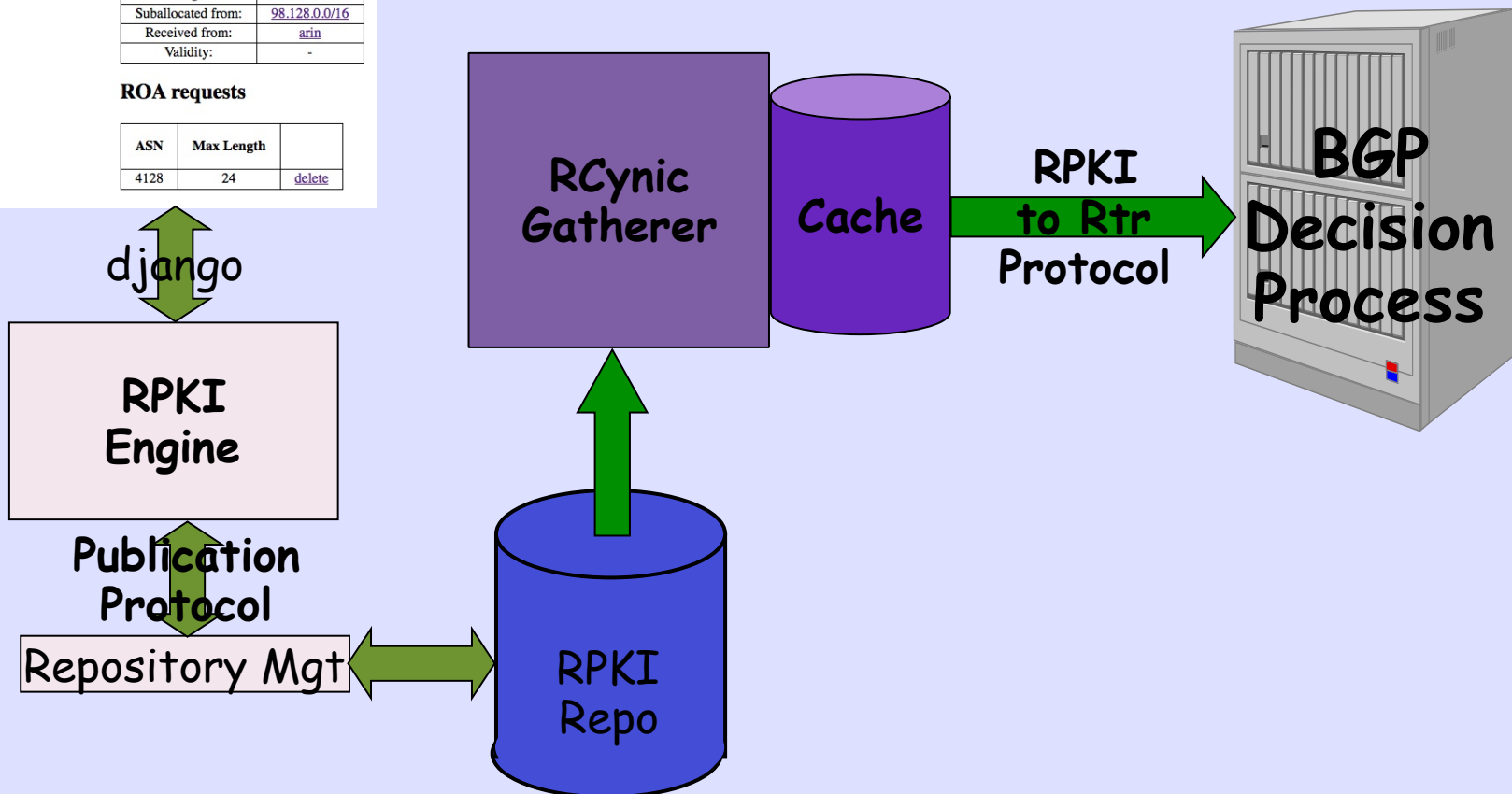
rgnet > Prefix View > 98.128.0.0/24

### Prefix View

Range:	98.128.0.0/24
Suballocated from:	98.128.0.0/16
Received from:	arin
Validity:	-

### ROA requests

ASN	Max Length	
4128	24	<a href="#">delete</a>



# Assume Origin Validation

```
R3#sh ip bg 98.128.0.0/24
```

```
BGP routing table entry for 98.128.0.0/24, version 94
```

```
Paths: (2 available, best #2, table default)
```

```
65000 3130
```

```
10.0.0.1 from 10.0.0.1 (65.38.193.12)
```

```
Origin IGP, localpref 100, valid, external  
path 6802D4DC RPKI State invalid
```

```
65001 4128
```

```
10.0.1.1 from 10.0.1.1 (65.38.193.13)
```

```
Origin IGP, localpref 100, valid, external, best  
path 6802D7C8 RPKI State valid
```

# Origin Validation is Weak

- Today's Origin Validation provides neither cryptographic assurance (announcements are not signed), nor assurance of the AS Path of the announcement.
- A malicious router may announce as any AS, i.e. forge the ROAed origin AS.
- This would pass ROA Validation

# Protocol Not Policy

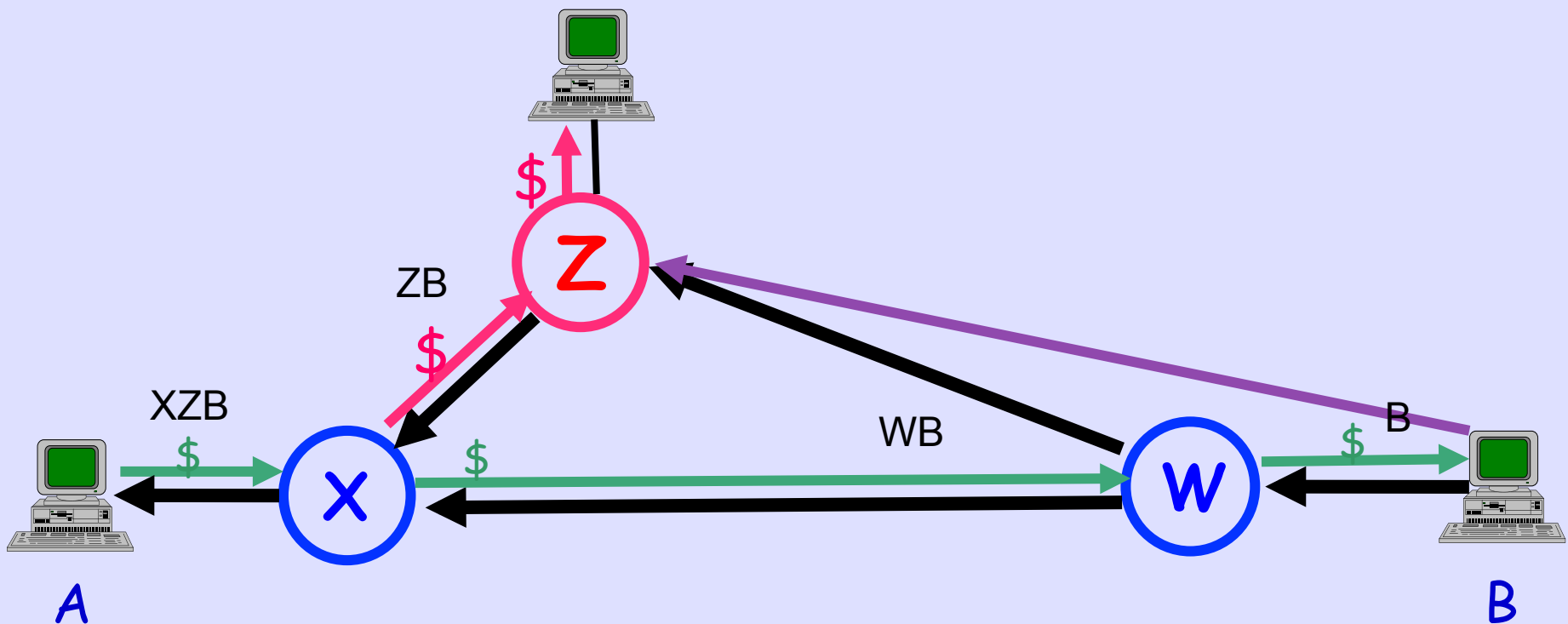
- Policy on the global Internet changes every 36ms
- We already have a protocol to distribute policy or its effects, it is called BGP
- We can not know intent, should Mary have announced the prefix to Bob
- But Joe can formally validate that Mary did announce the prefix to Bob
- BGPsec validates that the protocol has not been violated, and is not about intent or business policy

# Full Path Validation

- Rigorous per-prefix AS path validation is the goal
- Protect against origin forgery and AS-Path monkey in the middle attacks
- Not merely showing that a received AS path is not impossible
- Yes, this is S-BGP-like not SO-BGP-like



# Path Shortening Attack

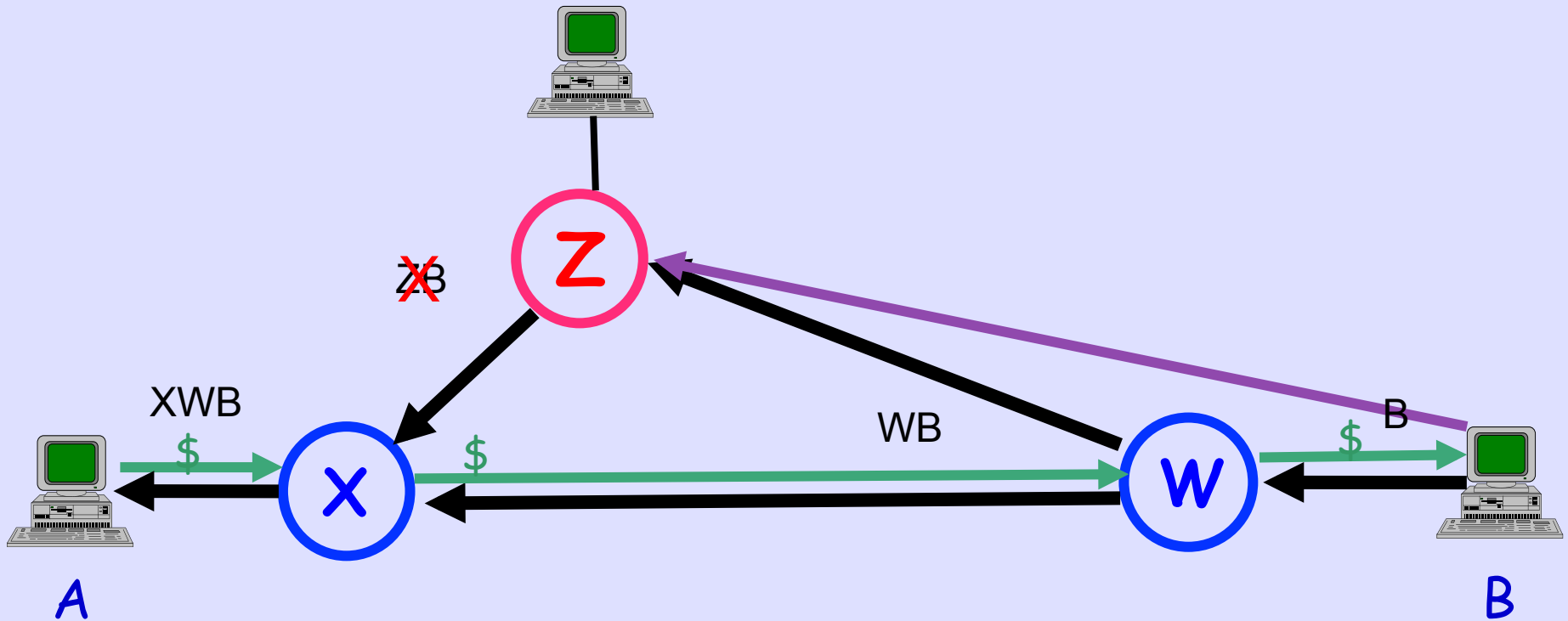


Expected Path - A->X->W->B

Diverted Path - A->X->Z->W->B

There Are Many Many Other Attacks

# Forward Path Signing

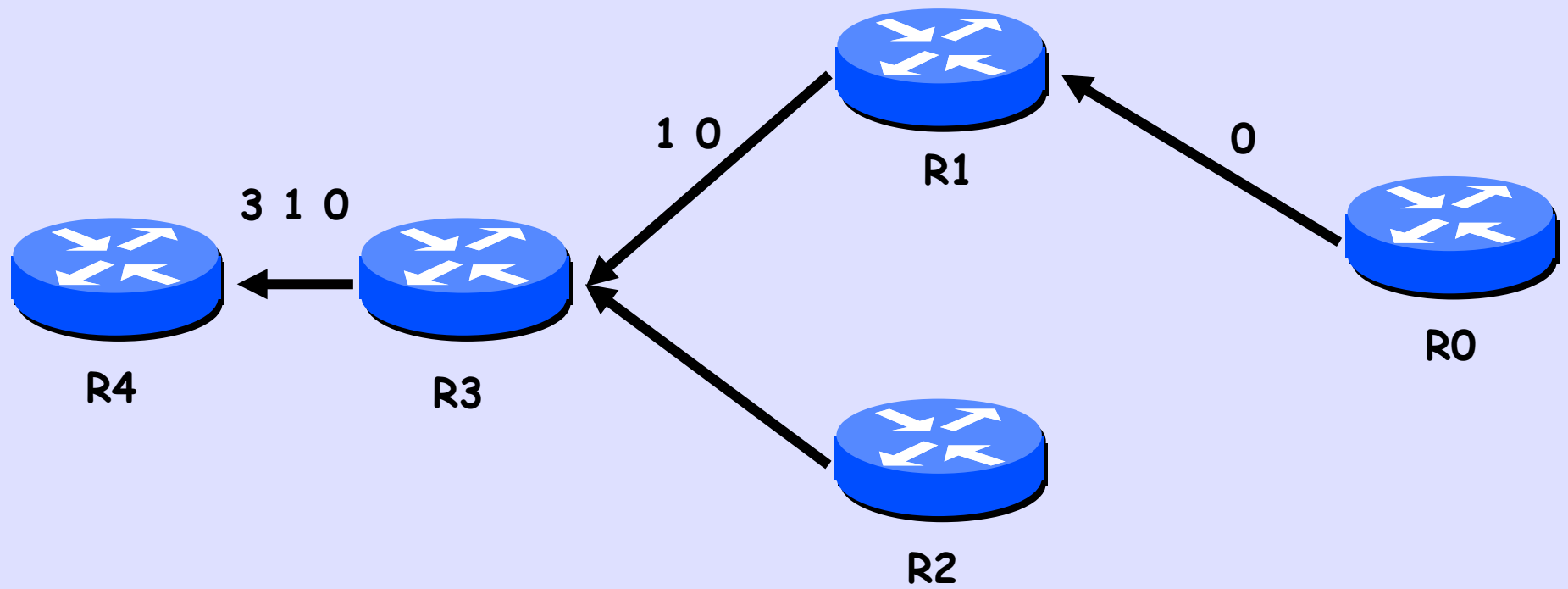


B cryptographically signs the message to W  $S_b(B \rightarrow W)$   
W signs messages to X and Z encapsulating B's message  
 $S_w(W \rightarrow X (S_b(B \rightarrow W)))$  and  $S_w(W \rightarrow Z (S_b(B \rightarrow W)))$   
X signs the message to A  $S_x(X \rightarrow A (S_w(W \rightarrow X (S_b(B \rightarrow W)))))$   
Z can only sign  $S_z(Z \rightarrow X (S_w(W \rightarrow Z (S_b(B \rightarrow W)))))$

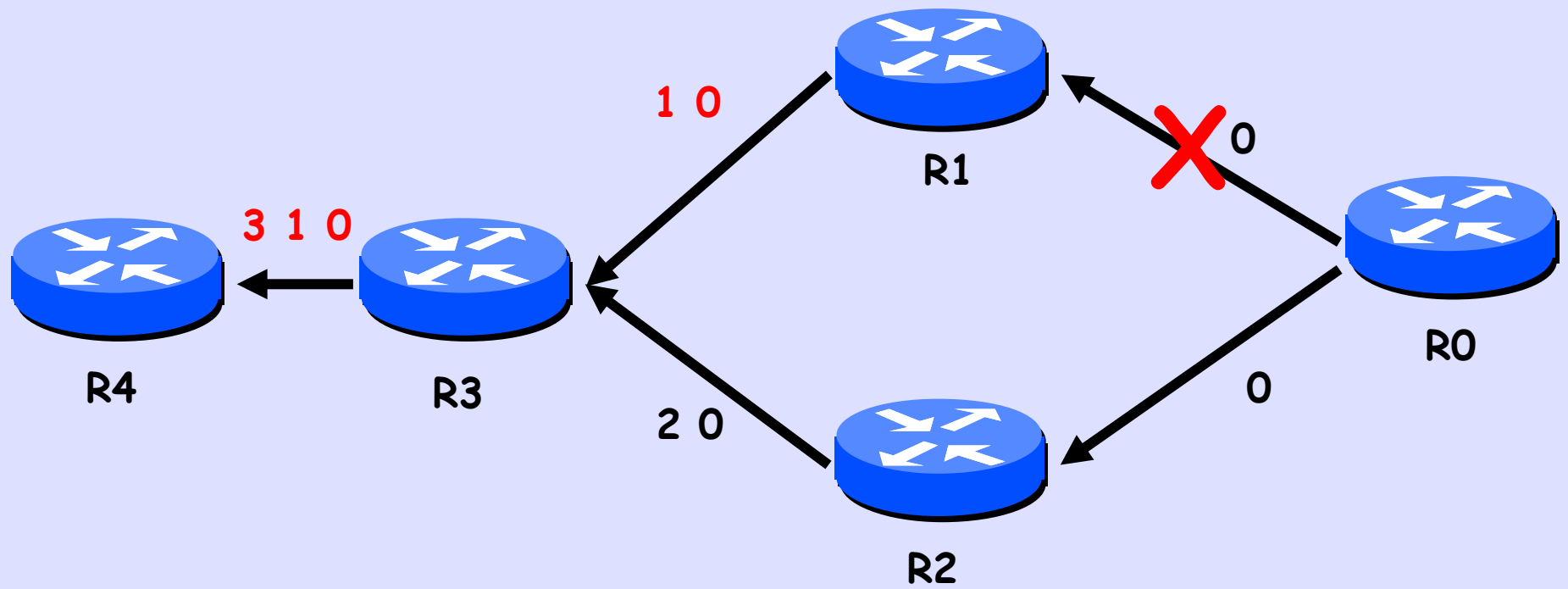
# Capability Negotiation

- It is assumed that consenting routers will use BGP capability exchange to agree to run BGPsec between them
- The capability will, among other things remove the 4096 PDU limit for updates
- If BGPsec capability is not agreed, then only traditional BGP data are sent

# Replay Attack



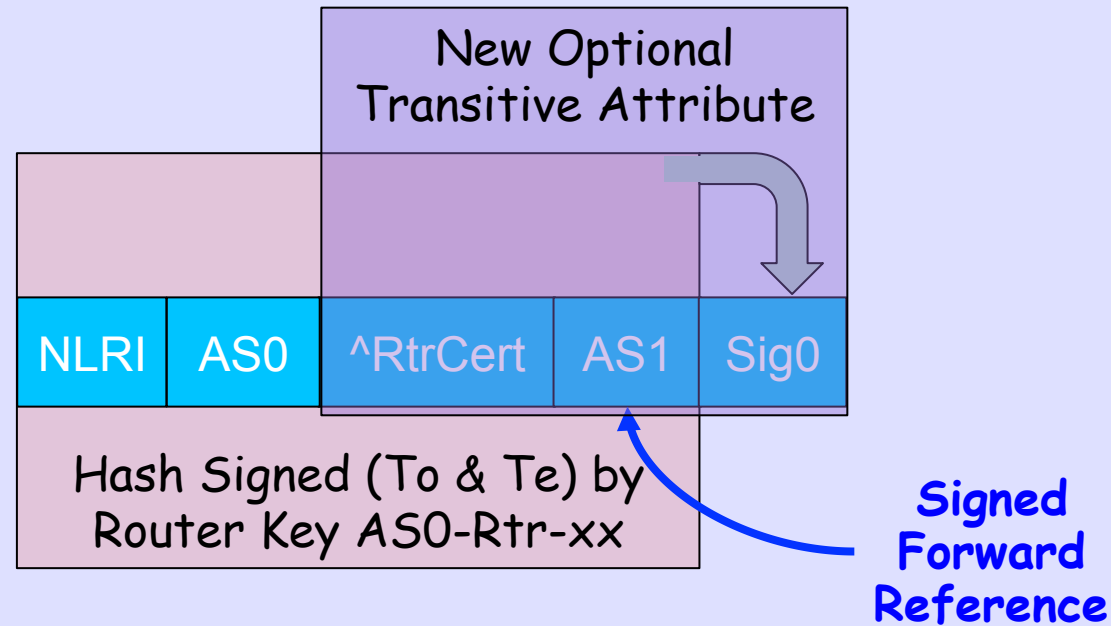
# Replay Attack



# Replay Reduction

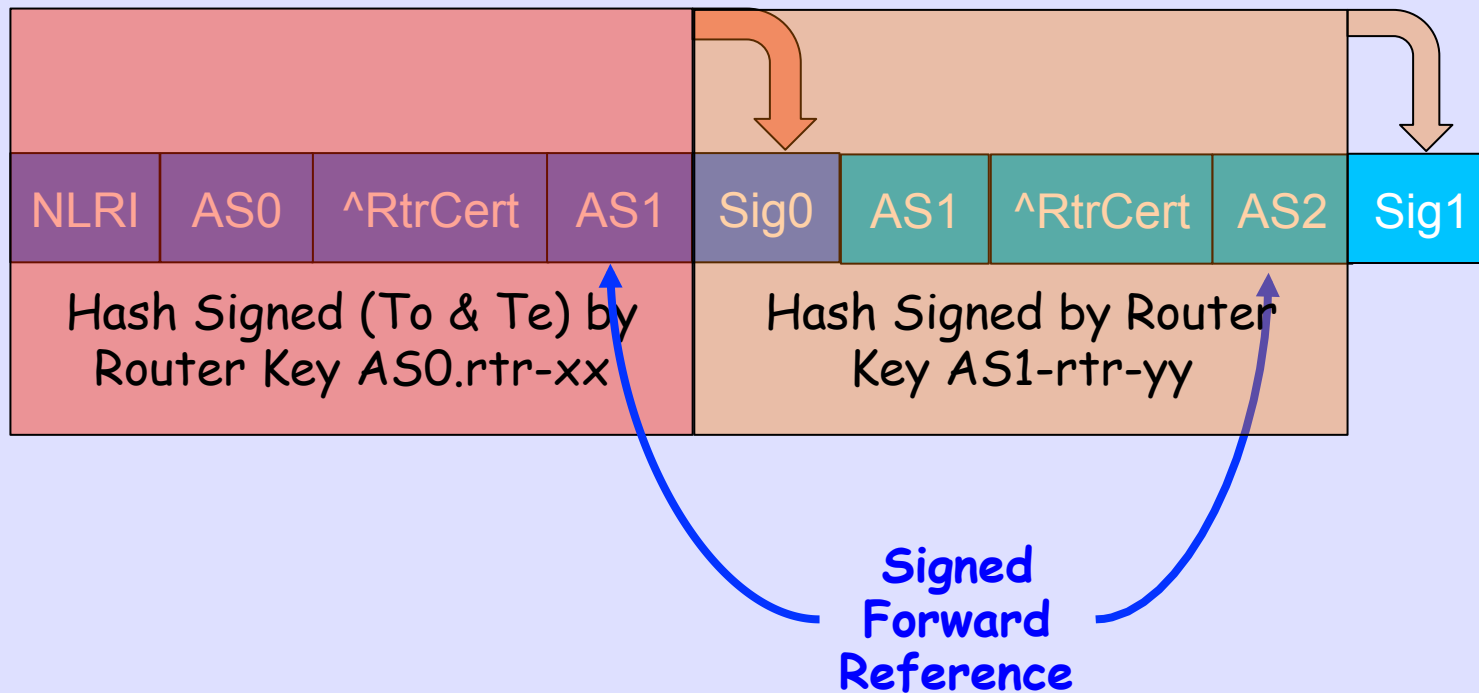
- Announcement replay is a vulnerability
- Therefore freshness is critical
- So originating announcer signs with a relatively short signature lifetime
- Origin re-announces prefix well within that lifetime, *AKA beaconing*
- Suggested to be days, but can be hours for truly critical infrastructure

# Origination by AS0 to AS1



- To and Te are times of signature origination and expiration
- Signature has a well-jittered validity end time, Te, of days
- Re-announcement by origin, AKA *beaconing*, every  $\sim (Te - To) / 3$
- ROA is not needed as prefix is sufficient to find it in RPKI as today

# Announcement AS1 to AS2



- R1 signing over R0's signature is same as signing over entire R0 announcement
- Non-originating router signatures do not have validity periods
- But when they receive a beacon announcement, they must propagate it



# Only at Provider Edges

- This design protects only inter-domain routing, not IGPs, not even iBGP
- BGPsec will be used inter-provider, only at the providers' edges
- Of course, the provider's iBGP will have to carry the BGPsec information
- Providers and inter-provider peerings might be heterogeneous

# Simplex End Site

- End Site can trust up-stream's policies
- But they want their origination signed
- So they announce capability to send but do not accept signed data
- They sign announcement and beacon
- This can be done without hardware upgrade!!

# Some Consequences



# New Hardware Generation

It is likely that routers will have to be upgraded to use this design, likely with much more memory and probably with hardware crypto assistance. It is accepted that this means that it will be some years,  $O(\text{IPv6 ASIC upgrades})$  before there is more than test deployment

# Route Servers

BGPsec can't forward sign across an AS-transparent route server as you do not know the peer AS

# Proxy Aggregation

Proxy Aggregation, i.e.

AS-Sets, is not supported

# Does Not Lock Data Plane

- It is acknowledged that rigorous control plane verification does not in any way guarantee that packets follow the control plane
- See IMC 2009 paper which shows that 70% of the ASs in the so-called 'default free zone' also have default

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**we take your scissors away and turn them into plowshares**